



Interval Methods for Reliable Modeling, Identification and Control of Dynamic Systems

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Motivation Arithmetics System Formulations Problem: Overestimation and Ho

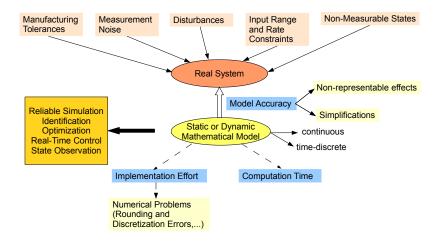
Part 1

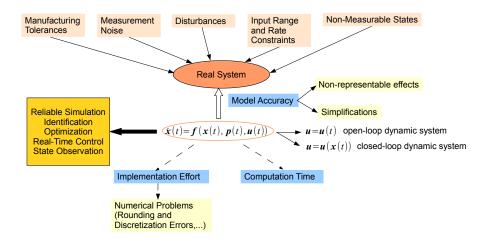
1. Fundamentals of Interval Arithmetic

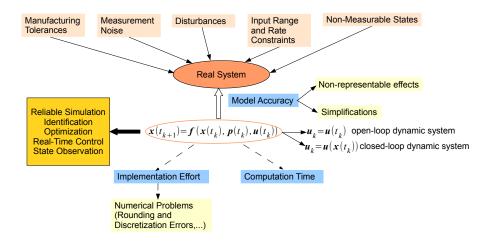
Presentation of the Fundamental Mathematical Concept of Interval Arithmetic for Set-Valued Computations Presentation of the Fundamental Mathematical Concept of Interval Arithmetic for Set-Valued Computations

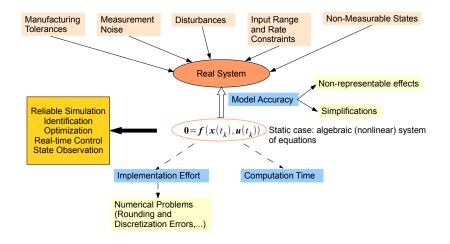
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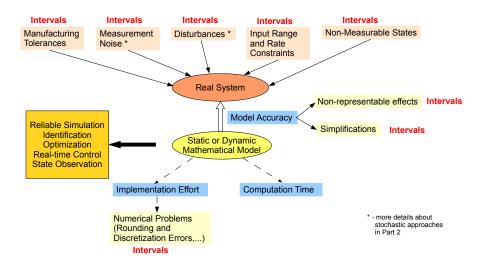
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Definition of Real Intervals, Interval Vectors, Interval Matrices

Scalar Real Interval

 $[a] = [\underline{a}; \overline{a}] = [\inf([a]); \sup([a]))] \ , \ \underline{a} \leq \overline{a} \ , \ \{x \in \mathbb{R} | \underline{a} \leq x \leq \overline{a}\}$

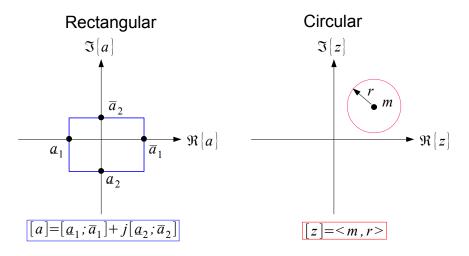
Interval Vector

$$[\mathbf{a}] = \begin{bmatrix} [\underline{a}_1; \overline{a}_1] \\ [\underline{a}_2; \overline{a}_2] \\ \vdots \\ [\underline{a}_n; \overline{a}_n] \end{bmatrix}$$

Interval Matrix

$$[\mathbf{A}] = \begin{bmatrix} \underline{[a_{11}; \overline{a}_{11}]} & \underline{[a_{12}; \overline{a}_{12}]} & \dots & \underline{[a_{1n}; \overline{a}_{1n}]} \\ \underline{[a_{21}; \overline{a}_{21}]} & \underline{[a_{22}; \overline{a}_{22}]} & \dots & \underline{[a_{2n}; \overline{a}_{2n}]} \\ \vdots & \vdots & \ddots & \vdots \\ \underline{[a_{n1}; \overline{a}_{n1}]} & \underline{[a_{n2}; \overline{a}_{n2}]} & \dots & \underline{[a_{nn}; \overline{a}_{nn}]} \end{bmatrix}$$

Definition of Complex Intervals



 \Rightarrow Useful for dynamic systems with oscillatory behavior

Addition

$$[\underline{p};\overline{p}] + [\underline{q};\overline{q}] = [\underline{p} + \underline{q};\overline{p} + \overline{q}]$$

$$[1;2] + [-2;2] = [1 + (-2);2 + 2] = [-1;4]$$

$$\begin{bmatrix} [-2;-1]\\[0;4] \end{bmatrix} + \begin{bmatrix} [-10;-3]\\[5;8] \end{bmatrix} = \begin{bmatrix} [-12;-4]\\[5;12] \end{bmatrix}$$

$$\begin{bmatrix} [2;3] & [-4;-3] \\ [7;9] & [10;15] \end{bmatrix} + \begin{bmatrix} [12;13] & [-14;-13] \\ [17;19] & [20;25] \end{bmatrix} = \begin{bmatrix} [14;16] & [-18;-16] \\ [24;28] & [30;40] \end{bmatrix}$$

Subtraction

$$[\underline{p};\overline{p}] - [\underline{q};\overline{q}] = [\underline{p} - \overline{q};\overline{p} - \underline{q}]$$

$$[1;2] - [2;3] = [1 - 3;2 - 2] = [-2;0]$$

$$\begin{bmatrix} [2;3] & [-4;-3] \\ [7;9] & [10;15] \end{bmatrix} - \begin{bmatrix} [12;13] & [-14;-13] \\ [17;19] & [20;25] \end{bmatrix} = \begin{bmatrix} [-11;-9] & [9;11] \\ [-12;-8] & [-15;-5] \end{bmatrix}$$

Multiplication

 $[\underline{p};\overline{p}] \cdot [\underline{q};\overline{q}] = \left[\min\left\{\underline{p}\,\underline{q},\underline{p}\,\overline{q},\overline{p}\,\underline{q},\overline{p}\,\overline{q}\right\}; \max\left\{\underline{p}\,\underline{q},\underline{p}\,\overline{q},\overline{p}\,\underline{q},\overline{p}\,\overline{q}\right\}\right]$

 $[1;2] \cdot [2;3] = [\min\{1 \cdot 2, 1 \cdot 3, 2 \cdot 2, 2 \cdot 3\}; \max\{1 \cdot 2, 1 \cdot 3, 2 \cdot 2, 2 \cdot 3\}] = [2;6]$

Division

$$\tfrac{[p]}{[q]} = [p] \cdot \left[\tfrac{1}{\overline{q}} \ ; \ \tfrac{1}{\underline{q}} \right] \quad \text{ if } \quad 0 \not \in [q]$$

$$\frac{[1;2]}{[2;3]} = [1;2] \cdot \left[\frac{1}{3} \ ; \ \frac{1}{2}\right] = \left[\frac{1}{3};1\right]$$

Radius of a Real Interval

$$r([a]) = \frac{1}{2}(\overline{a} - \underline{a})$$

Width of an Interval

$$w([a]) = \overline{a} - \underline{a} = 2 \cdot r([a])$$

Mid-point of an Interval

$$m([a]) = \frac{1}{2}(\underline{a} + \overline{a})$$

 \Rightarrow For real interval vectors and matrices, these characteristics hold component-wise

Continuous- and Discrete-Time Systems — Dynamic Case

Continuous-Time System $\dot{\mathbf{x}}(t) = \mathbf{f}(\mathbf{x}(t), \mathbf{p}, \mathbf{u}(t))$

 $\mathbf{x}(t)$ State Vector

p Vector of Uncertain Parameters: $p_i \in \left[\underline{p}_i \ ; \ \overline{p}_i\right], i = 1, ..., n_p$

 $\mathbf{u}(t)$ Input Vector: $u_j \in \left[\underline{u}_j \ ; \ \overline{u}_j\right], j = 1, ..., n_u$

Continuous- and Discrete-Time Systems — Dynamic Case

Continuous-Time System $\dot{\mathbf{x}}(t) = \mathbf{f}(\mathbf{x}(t), \mathbf{p}, \mathbf{u}(t))$

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Discrete-Time System $\mathbf{x}(t_{k+1}) = \mathbf{f}(\mathbf{x}(t_k), \mathbf{p}(t_k), \mathbf{u}(t_k))$

- $\mathbf{x}(t_k)$ State Vector
- $\begin{array}{ll} \mathbf{p} & \mbox{Vector of Uncertain Parameters: } p_i \in \left\lfloor \underline{p}_i \ ; \ \overline{p}_i \right\rfloor \\ & \Longrightarrow \mbox{Range Bounds / Tolerances} \\ \mathbf{u}(t_k) & \mbox{Input Vector: } u_j(t_k) \in \left[\underline{u}_j(t_k) \ ; \ \overline{u}_j(t_k)\right] \\ & \Longrightarrow \mbox{Input Range Constraints} \end{array}$

 \Rightarrow Calculate all reachable states

Continuous- and Discrete-Time Systems — Static Case

Continuous-Time System $\dot{\mathbf{x}}(t) = \mathbf{0} = \mathbf{f}(\mathbf{x}(t), \mathbf{p}, \mathbf{u}(t))$

 $\mathbf{x}(t)$ State Vector

p Vector of Uncertain Parameters: $p_i \in \left| \underline{p}_i ; \overline{p}_i \right|, i = 1, ..., n_p$

 $\mathbf{u}(t)$ Input Vector: $u_j \in \left[\underline{u}_j \ ; \ \overline{u}_j\right], j = 1, ..., n_u$

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 \Rightarrow Solve for state vector $\mathbf{x}(t)$ or \mathbf{x}_k resp. for 1 time step

Problem: Multiple occurence of an interval in one equation Necessary: Factorizations, simplifications, reformulations as far as possible

 \Rightarrow Reduction of overestimation and computation time

Example

$$[f]([x]) = 2 \cdot [x] - [x] \cdot [x]$$
 and $[x] = [-1 \ ; \ 2]$, Results provided by Intlab

1
$$[f]([x]) = 2 \cdot [x] - [x] \cdot [x] = [-6; 6]$$

2
$$[f]([x]) = 2 \cdot [x] - [x]^2 = [-6; 4]$$

3
$$[f]([x]) = -([x] - 1) \cdot ([x] - 1) + 1 = [-3; 3]$$

• $[f]([x]) = -([x] - 1)^2 + 1 = [-3; 1]$ (Exact evaluation)

Problem: Multiple occurence of an interval in one equation

Necessary: Factorizations, simplifications, reformulations as far as possible \Rightarrow Reduction of overestimation and computation time

Example

 $[f]([x]) = 2 \cdot [x] - [x] \cdot [x]$ and [x] = [-1; 2], Results provided by Intlab Higher-order Interval Evaluation for Polynomials: Taylor Expansion $(x_m = mid([x]))$ according to

$$T(f) = f(x_m) + \left(\sum_{i=1}^{n-1} \left. \frac{\partial^i f(x)}{\partial x^i} \right|_{x_m} \cdot \frac{([x] - x_m)^i}{i!} \right) + \left. \frac{\partial^n f(x)}{\partial x^n} \right|_{[x]} \cdot \frac{([x] - x_m)^n}{n!}$$

Problem: Multiple occurence of an interval in one equation

Necessary: Factorizations, simplifications, reformulations as far as possible \Rightarrow Reduction of overestimation and computation time

Example

$$[f]([x]) = 2 \cdot [x] - [x] \cdot [x]$$
 and $[x] = [-1; 2]$, Results provided by Intlab

$$\begin{aligned} \text{Faylor Expansion with } x_m &= \operatorname{mid}([x]) = 0.5\\ [f]([x]) &= f(x_m) + \left. \frac{\partial f}{\partial x} \right|_{x_m} \cdot ([x] - x_m) + \left. \frac{\partial^2 f}{\partial x^2} \right|_{[x]} \cdot \frac{([x] - x_m)^2}{2!}\\ f(x_m) &= 2 \cdot 0.5 - 0.5^2 = 0.75\\ \left. \frac{\partial f}{\partial x} \right|_{x_m} \cdot ([x] - x_m) = (2 - 2 \cdot x)|_{x_m} \cdot ([x] - x_m) = [-1.5 \ ; \ 1.5] \end{aligned}$$

Example

$$[f]([x])=2\cdot [x]-[x]\cdot [x]$$
 and $[x]=[-1\ ;\ 2],$ Results provided by Intlab

• Taylor Expansion with
$$x_m = \operatorname{mid}([x]) = 0.5$$

 $[f]([x]) = 0.75 + [-1.5; 1.5] + \frac{\partial^2 f}{\partial x^2}\Big|_{[x]} \cdot \frac{([x] - x_m)^2}{2!} = \frac{\partial^2 f}{\partial x^2}\Big|_{[x]} \cdot \frac{([x] - x_m) \cdot ([x] - x_m)}{2!} = -2 \cdot \frac{([x] - x_m) \cdot ([x] - x_m)}{2!} = -1 \cdot [-1.5; 1.5] \cdot [-1.5; 1.5] = [-2.25; 2.25]$

 \Rightarrow Taylor expansion will be demonstrated later with two software libraries

Example

 $[f]([x]) = 2 \cdot [x] - [x] \cdot [x]$ and [x] = [-1; 2], Results provided by Intlab **(**) Taylor Expansion with $x_m = mid([x]) = 0.5$ $[f]([x]) = 0.75 + [-1.5; 1.5] + \frac{\partial^2 f}{\partial x^2}\Big|_{[x]} \cdot \frac{([x] - x_m)^2}{2!} =$ $\frac{\partial^2 f}{\partial x^2}\Big|_{[x]} \cdot \frac{([x] - x_m) \cdot ([x] - x_m)}{2!} = -2 \cdot \frac{([x] - x_m) \cdot ([x] - x_m)}{2!} =$ $-1 \cdot [-1.5; 1.5] \cdot [-1.5; 1.5] = [-2.25; 2.25]$ $[x]^{2} = \begin{cases} [\min(\underline{aa}, \overline{aa}) \ ; \ \max(\underline{aa}, \overline{aa})] & \text{if } 0 \notin [x] \\ [0 \ ; \ \max(\underline{aa}, \overline{aa})] & \text{if } 0 \in [x] \end{cases}$ $\frac{\partial^2 f}{\partial x^2}\Big|_{[x]} \cdot \frac{([x] - x_m)^2}{2!} = -2 \cdot \frac{([x] - x_m)^2}{2!} = -1 \cdot [-1.5 \ ; \ 1.5]^2 = [-2.25 \ ; \ 0]$

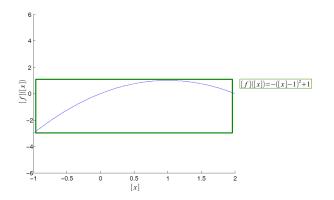
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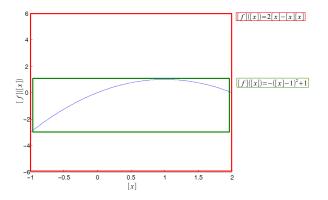
Example

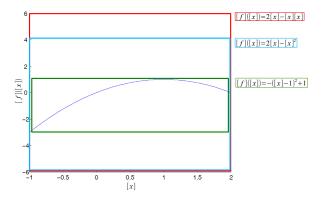
 $\begin{array}{l} [f]([x]) = 2 \cdot [x] - [x] \cdot [x] \text{ and } [x] = [-1 \ ; \ 2], \text{ Results provided by Intlab} \\ \textcircled{S} \quad \text{Taylor Expansion with } x_m = \operatorname{mid}([x]) = 0.5 \\ [f]([x]) = f(x_m) + \left. \frac{\partial f}{\partial x} \right|_{x_m} \cdot ([x] - x_m) + \left. \frac{\partial^2 f}{\partial x^2} \right|_{[x]} \cdot \frac{([x] - x_m)^2}{2!} = \end{array}$

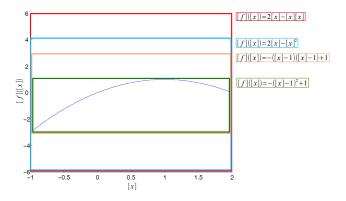
$$\begin{split} f(x_m) &= 2 \cdot 0.5 - 0.5^2 = 0.75\\ \frac{\partial f}{\partial x}\Big|_{x_m} \cdot ([x] - x_m) = (2 - 2 \cdot x)\Big|_{x_m} \cdot ([x] - x_m) = [-1.5 \ ; \ 1.5]\\ \frac{\partial^2 f}{\partial x^2}\Big|_{[x]} \cdot \frac{([x] - x_m)^2}{2!} &= -2 \cdot \frac{([x] - x_m)^2}{2!} = -1 \cdot [-1.5 \ ; \ 1.5]^2 = [-2.25 \ ; \ 0]\\ \Rightarrow [f]([x]) &= 0.75 + [-1.5 \ ; \ 1.5] + [-2.25 \ ; \ 0] = [-3 \ ; \ 2.25] \end{split}$$

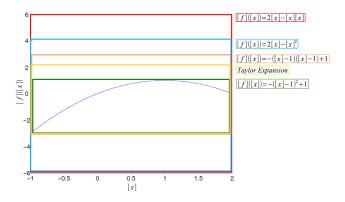
 \Rightarrow Taylor expansion will be demonstrated later with two software libraries



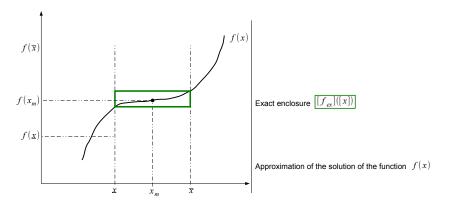




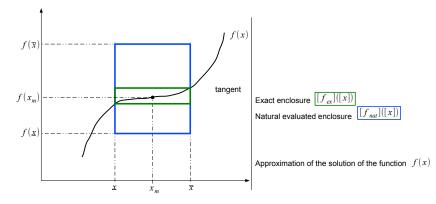




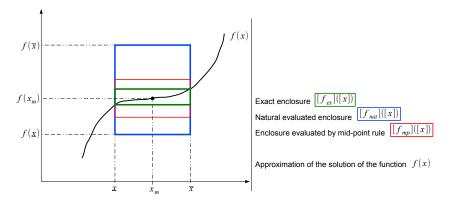
$$f(x) \subseteq f_m([x]) = f(x_m) + \left. \frac{\partial f}{\partial x} \right|_{[x]} ([x] - x_m)$$



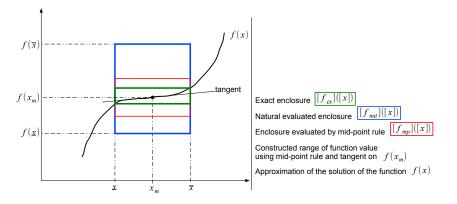
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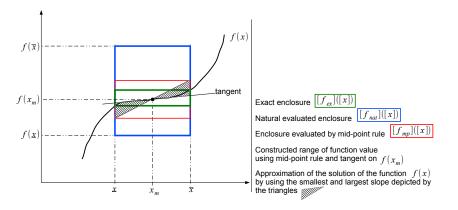
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$$f(x) \subseteq f_m([x]) = f(x_m) + \left. \frac{\partial f}{\partial x} \right|_{[x]} ([x] - x_m)$$



$$f(x) \subseteq f_m([x]) = f(x_m) + \left. \frac{\partial f}{\partial x} \right|_{[x]} ([x] - x_m)$$



Monotonicity

Consider: Interval-Valued Function given by $F = x + x \cdot x$

- Two intervals $[x_1] = [-2; 4]$ and $[x_2] = [-1; 4]$ with $[x_1] \subset [x_2]$
- $F([x_1]) = [-2;4] + [-2;4] \cdot [-2;4] = [-2;4] + [-8;16] = [-10;20]$
- $F([x_2]) = [-1;4] + [-1;4] \cdot [-1;4] = [-1;4] + [-4;16] = [-5;20]$
- Consequence $F([x_2]) \subset F([x_1]) \Rightarrow F$ is an inclusion monotonic function
- 4 basic arithmetic operators are also inclusion monotonic

Consequence for Calculating with Intervals

- Splitting of large intervals
- Hull of all evaluations with the subintervals
- Tighter range bounds than with original interval

Monotonicity

Consider: Interval-Valued Function given by $F = x + x \cdot x$

- Two intervals $[x_1] = [-2;4]$ and $[x_2] = [-1;4]$ with $[x_1] \subset [x_2]$
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- Consequence $F([x_2]) \subset F([x_1]) \Rightarrow F$ is an inclusion monotonic function
- 4 basic arithmetic operators are also inclusion monotonic

Monitonicity of a Function Using Derivatives

$$\left. \frac{\partial F}{\partial \mathbf{x}} \right|_{x \in [x]} < 0 \quad \Rightarrow F \in [F(\overline{x}); F(\underline{x})]$$

$$\left. \frac{\partial F}{\partial \mathbf{x}} \right|_{x \in [x]} > 0 \quad \Rightarrow F \in [F(\underline{x}); F(\overline{x})]$$

Overestimation: Wrapping Effect —Example

Discrete System Model

$$[\mathbf{x}](t_{k+1}) = \mathbf{A} \cdot [\mathbf{x}](t_k) \text{ with } [\mathbf{x}](t_0) = \begin{bmatrix} [-1 \ ; \ 1] \\ [-1 \ ; \ 1] \end{bmatrix} \text{ and } \mathbf{A} = \frac{1}{2}\sqrt{2} \begin{bmatrix} 1 & 1 \\ -1 & 1 \end{bmatrix}$$

Aim

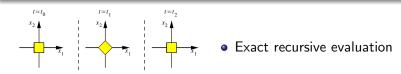
Evaluation of interval enclosure $[\mathbf{x}](t_{k+1})$

Problem in Engineering Tasks

Uncertainty in parameters, significantly larger than representation errors of floating-point values (rounding errors)

Overestimation: Wrapping Effect — Example

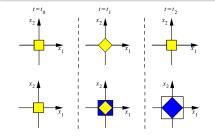
$$[\mathbf{x}](t_{k+1}) = \mathbf{A} \cdot [\mathbf{x}](t_k) \text{ with } [\mathbf{x}](t_0) = \begin{bmatrix} [-1 \ ; \ 1] \\ [-1 \ ; \ 1] \end{bmatrix} \text{ and } \mathbf{A} = \frac{1}{2}\sqrt{2} \begin{bmatrix} 1 & 1 \\ -1 & 1 \end{bmatrix}$$



 \Rightarrow Rotation of 45° due to structure of system matrix ${\bf A}$

Overestimation: Wrapping Effect — Example

$$[\mathbf{x}](t_{k+1}) = \mathbf{A} \cdot [\mathbf{x}](t_k), \ [\mathbf{x}](t_0) = \begin{bmatrix} [-1 \ ; \ 1] \\ [-1 \ ; \ 1] \end{bmatrix}, \ \mathbf{A} = \mathbf{A}_k = \frac{1}{2}\sqrt{2} \begin{bmatrix} 1 & 1 \\ -1 & 1 \end{bmatrix}$$



 Traditional recursive interval evaluation (using e.g. Intlab)

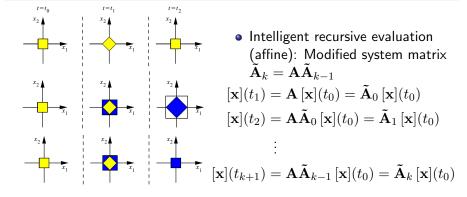
$$[\mathbf{x}](t_1) = \mathbf{A} [\mathbf{x}](t_0)$$
$$[\mathbf{x}](t_2) = \mathbf{A} [\mathbf{x}](t_1)$$
$$\vdots$$

$$[\mathbf{x}](t_{k+1}) = \mathbf{A} [\mathbf{x}](t_k)$$

 \Rightarrow Exponential growth of the enclosing interval boxes

Overestimation: Wrapping Effect — Example

$$[\mathbf{x}](t_{k+1}) = \mathbf{A} \cdot [\mathbf{x}](t_k), \ [\mathbf{x}](t_0) = \begin{bmatrix} [-1 \ ; \ 1] \\ [-1 \ ; \ 1] \end{bmatrix}, \ \mathbf{A} = \mathbf{A}_k = \frac{1}{2}\sqrt{2} \begin{bmatrix} 1 & 1 \\ -1 & 1 \end{bmatrix}$$



 \Rightarrow Significant reduction of the wrapping effect for linear systems

Affine System Representation for Discrete Systems

Advantages

- Directly mapping of interval variables to their initial intervals in each time step
- No dependencies between intervals \Rightarrow no interval box rotations

Discretization depends on

- Additive interval or multiplicatively coupled parameter interval
- Input variable constant or changing
- Explicit or implicit Euler method

Affine System Representation for a SISO System Case 1: Additive interval uncertainty [a], $u(t_k) \neq \text{const}$, step size T = 1 $f(y(t_k), [u(t_k)]) = 2 \cdot [y](t_k) + 1 \cdot u(t_k) + 3 \cdot [a]$ $[\mathbf{x}](t_{k+1}) = \underbrace{\begin{bmatrix} [y](t_{k+1}) \\ [a](t_{k+1}) \end{bmatrix}}_{\text{extended state vector } \mathbf{x}(t_{k+1})} = \underbrace{\begin{bmatrix} 2 & 3 \\ 0 & 1 \end{bmatrix}}_{\mathbf{M}} \cdot \underbrace{\begin{bmatrix} [y](t_k) \\ [a](t_k) \end{bmatrix}}_{\mathbf{x}(t_k)} + \underbrace{\begin{bmatrix} 1 \\ 0 \end{bmatrix} \cdot u(t_k)}_{\rho(t_k)}$

E.g. Explicit Euler Discetization (time discretization error neglected) $[\mathbf{x}](t_{k+1}) = \mathbf{M}(t_{k+1}) \cdot [\mathbf{x}](t_0) + \boldsymbol{\gamma}(t_{k+1}) \quad \text{with}$

$$\begin{split} \mathbf{M}(t_k) &= \mathbf{M} \quad \Rightarrow \quad \mathbf{M}(t_{k+1}) = \mathbf{M} \cdot \mathbf{M}(t_k) \\ \mathbf{\gamma}(t_{k+1}) &= \mathbf{M}(t_k) \cdot \mathbf{\gamma}(t_k) + T \cdot \mathbf{\rho}(t_k) \\ \text{initial conditions } \mathbf{x}(t_0) &= \begin{bmatrix} [y](t_0) \\ [a](t_0) \end{bmatrix}, \quad \mathbf{M}(t_0) = \mathbf{I}^{2 \times 2}, \quad \mathbf{\gamma}(t_0) = \mathbf{0} \end{split}$$

Affine System Representation for a SISO System Case 2: Additive interval uncertainty [a], u=const, step size T = 1

$$\begin{split} f(y(t), [u(t)]) &= 2 \cdot [y](t) + 1 \cdot u(t) + 3 \cdot [a] \\ [\mathbf{x}](t_{k+1}) &= \underbrace{\begin{bmatrix} [y](t_{k+1}) \\ [a](t_{k+1}) \\ u(t_{k+1}) \end{bmatrix}}_{\text{extended state vector } \mathbf{x}(t_{k+1})} = \underbrace{\begin{bmatrix} 2 & 3 & 1 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}}_{\mathbf{M}} \cdot \underbrace{\begin{bmatrix} [y](t_k) \\ [a](t_k) \\ u(t_k) \end{bmatrix}}_{\mathbf{x}(t_k)} \end{split}$$

Explicit Euler Discetization (time discretization error neglected)

$$[\mathbf{x}](t_{k+1}) = \mathbf{M}(t_{k+1}) \cdot [\mathbf{x}](t_0)$$
 with

$$\mathbf{M}(t_k) = \mathbf{M} \quad \Rightarrow \quad \mathbf{M}(t_{k+1}) = \mathbf{M} \cdot \mathbf{M}(t_k)$$
$$[\mathbf{x}](t_0) = \begin{bmatrix} [y](t_0) \\ [a](t_0) \\ u(t_0) \end{bmatrix}$$

Affine System Representation — Comparison

Implicit Euler Method $(u(t_{k+1}) = u(t_k) = const)$

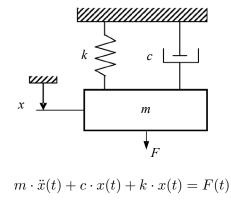
$$\begin{split} \dot{\mathbf{x}}(t) &= \mathbf{f}(\mathbf{x}(t), u(t)) = \mathbf{A}\mathbf{x}(t) + \mathbf{b}u(t) \\ \Rightarrow &\mathbf{f}(\mathbf{x}(t_{k+1}), u(t_{k+1})) \approx \frac{\mathbf{x}(t_{k+1}) - \mathbf{x}(t_k)}{T} \\ \mathbf{x}(t_{k+1}) &= (\mathbf{I} - T \cdot \mathbf{A})^{-1} \cdot (\mathbf{x}(t_k) + T \cdot \mathbf{b} \cdot u(t_{k+1})) \\ \underbrace{\begin{bmatrix} \mathbf{x}(t_{k+1}) \\ u(t_{k+1}) \end{bmatrix}}_{\tilde{\mathbf{x}}(t_{k+1})} &= \underbrace{\begin{bmatrix} (\mathbf{I} - T \cdot \mathbf{A})^{-1} & (\mathbf{I} - T \cdot \mathbf{A})^{-1} \cdot T \cdot \mathbf{b} \\ \mathbf{0}^T & \mathbf{1} \end{bmatrix}}_{\tilde{\mathbf{A}}} \cdot \underbrace{\begin{bmatrix} \mathbf{x}(t_k) \\ u(t_k) \end{bmatrix}}_{\tilde{\mathbf{x}}(t_k)} \end{split}$$

Affine System Representation — Comparison

Explicit Euler Method $(u(t_{k+1}) = u(t_k) = \text{const})$

$$\begin{aligned} \dot{\mathbf{x}}(t) &= \mathbf{f}(\mathbf{x}(t), u(t)) = \mathbf{A}\mathbf{x}(t) + \mathbf{b}u(t) \\ \Rightarrow &\mathbf{f}(\mathbf{x}(t_k), u(t_k)) \approx \frac{\mathbf{x}(t_{k+1}) - \mathbf{x}(t_k)}{T} \\ \mathbf{x}(t_{k+1}) &= (\mathbf{I} + T \cdot \mathbf{A}) \cdot \mathbf{x}(t_k) + T \cdot \mathbf{b} \cdot u(t_k) \\ \underbrace{\begin{bmatrix} \mathbf{x}(t_{k+1}) \\ u(t_{k+1}) \end{bmatrix}}_{\tilde{\mathbf{x}}(t_{k+1})} = \underbrace{\begin{bmatrix} (\mathbf{I} + T \cdot \mathbf{A}) & T \cdot \mathbf{b} \\ \mathbf{0}^T & 1 \end{bmatrix}}_{\tilde{\mathbf{A}}} \cdot \underbrace{\begin{bmatrix} \mathbf{x}(t_k) \\ u(t_k) \end{bmatrix}}_{\tilde{\mathbf{x}}(t_k)} \end{aligned}$$

Affine System Representation for a One-Mass Oscillator



$$\begin{bmatrix} \dot{x}(t) \\ \ddot{x}(t) \end{bmatrix} = \begin{bmatrix} 0 & 1 \\ -\frac{k}{m} & -\frac{c}{m} \end{bmatrix} \begin{bmatrix} x(t) \\ \dot{x}(t) \end{bmatrix} + \begin{bmatrix} 0 \\ \frac{1}{m} \end{bmatrix} F(t)$$

Affine System Representation for a One-Mass Oscillator

$$m \cdot \ddot{x}(t) + c \cdot x(t) + k \cdot x(t) = F(t)$$

$$\begin{bmatrix} \dot{x}(t) \\ \ddot{x}(t) \end{bmatrix} = \begin{bmatrix} 0 & 0 \\ -\frac{k}{m} & -\frac{c}{m} \end{bmatrix} \begin{bmatrix} x(t) \\ \dot{x}(t) \end{bmatrix} + \begin{bmatrix} 0 \\ \frac{1}{m} \end{bmatrix} F(t)$$

Euler Discretization with constant input variable $\tilde{\mathbf{x}}(t_{k+1}) = \tilde{\mathbf{A}} \cdot \tilde{\mathbf{x}}(t_k)$

$$\begin{split} \tilde{\mathbf{x}}(t_1) &= \tilde{\mathbf{A}} \cdot \tilde{\mathbf{x}}(t_0) \\ \tilde{\mathbf{x}}(t_2) &= \tilde{\mathbf{A}} \cdot \tilde{\mathbf{x}}(t_1) = \tilde{\mathbf{A}} \cdot (\tilde{\mathbf{A}} \cdot \tilde{\mathbf{x}}(t_0)) = \tilde{\mathbf{A}}^2 \cdot \tilde{\mathbf{x}}(t_0) \\ \vdots \\ \tilde{\mathbf{x}}(t_{k+1}) &= \tilde{\mathbf{A}}^k \cdot \tilde{\mathbf{x}}(t_0) \end{split}$$

Affine System Representation for a One-Mass Oscillator

$$m \cdot \ddot{x}(t) + c \cdot x(t) + k \cdot x(t) = F(t)$$

$$\begin{bmatrix} \dot{x}(t) \\ \ddot{x}(t) \end{bmatrix} = \begin{bmatrix} 0 & 0 \\ -\frac{k}{m} & -\frac{c}{m} \end{bmatrix} \begin{bmatrix} x(t) \\ \dot{x}(t) \end{bmatrix} + \begin{bmatrix} 0 \\ \frac{1}{m} \end{bmatrix} F(t)$$

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Example in MATLAB

Motivation

Arithmetics

Part 1

1. Fundamentals of Interval Arithmetic

Software Demonstration

Software Demonstration of Interval Arithmetics

- INTLAB: INTerval LABoratory Matlab toolbox for Reliable Computing
- C-XSC C++ Class Library

Importance of Verified Computing

- Floating-point arithmetics on today's computer is always affected by a maximum accuracy
 ⇒ rounded results differ at most by 1 unit in the last place from the exact result
- After further calculations, the result may be wrong because of rounding ⇒ Results have to be verified

INTerval LABoratory

Development by Prof. Dr. Siegfried M. Rump, Hamburg University of Technology

http://www.ti3.tu-harburg.de/rump/intlab/

Standard interval arithmetic

- Arithmetic operators +, -, ·, /
- Real and complex intervals

Automatic Differentiation

- Forward mode: forward substitution to find the derivatives
- Compute derivatives using the chain rule for composite functions
- Calculate an enclosure of the true derivative of an interval function

INTerval LABoratory

http://www.ti3.tu-harburg.de/rump/intlab/

Verified Functions for Linear Systems of Equations

- Solution of linear systems of equations in a verified way
- Computation of an enclosure of the solution hull
- Aim: produce a tight bound on the true solution

Rounding Mode

- Function setround(y): changes the rounding mode of the processor to the nearest (0), round down (-1), round up (1)
- Function getround outputs the current rounding mode

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Matlab Example: intlab_fundamentals.m and taylor_expansion.m

			Problem: Overestimation and How to Reduce		
C-XSC					
http://www2.math.uni-wuppertal.de/~xsc/xsc/cxsc.html					

Information

- C++ Class Library for Extended Scientific Computing
- Compatible to Windows, Linux, Mac Os

Data Types

- real, interval, complex, cinterval (complex interval)
- rvector, ivector, cvector, civector (complex interval vector)
- rmatrix, imatrix, cmatrix, cimatrix (complex interval matrix)

	Problem: Overestimation and How to Reduce

C-XSC

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Rounding Mode

- by-default: all operations are only one rounding away from the exact result
- Modes: long fix-point accumulator for dot product computations (default), pure floating point operations, DotK algorithm (based on so-called error free transformations)

C-XSC

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Data Types fundamentals.cpp

- real, interval, complex, cinterval (complex interval)
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		Problem: Overestimation and How to Reduce
FADBA	D++	

http://www.fadbad.com/fadbad.html#General_introduction

General

- Flexible Automatic Differentiation using templates and operator overloading in C++
- Implementing the forward, backward and Taylor methods utilizing C++ templates and operator overloading
- Differentiate a C++ function by replacing all occurrences of the original arithmetic type with the AD-template version
- Possible to generate high-order derivatives

FADBAD++

http://www.fadbad.com/fadbad.html#General_introduction

General taylor_expansion_FF.cpp and taylor_expansion_T.cpp

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Advantage of Using C++ instead of Intlab

Interface to rapid control prototyping environments is possible

Motivation		

Thank you for your attention!

All presentations, examples and selected publications will be available at http://www.com.uni-rostock.de/ecc15/ in the 1st week of August

User: ECC15 Password: intervals-are-fun