



OPTIMAL SWITCHING INSTANTS FOR THE CONTROL OF HYBRID SYSTEMS

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Generic case

$$\left[\begin{array}{ll} \max_{u(\cdot)} & J(u(\cdot)) = \int_0^T h(x(t), u(t))dt + g(x(T)) \quad (\text{cost function}) \\ \text{s. t.} & \dot{x} = f(x(t), u(t)), \quad 0 < t \leq T \quad (\text{dynamical constraint}) \\ & x(0) = x_0, \quad h(x(t)) \in \mathcal{H}, \quad 0 < t \leq T \quad (\text{boundary conditions}) \\ & u(t) \in \mathcal{U}, \quad \forall t \quad (\text{bounded control}) \end{array} \right]$$

- The **dynamical constraint** coupled with the **boundary conditions** is an IVP-ODE depending on a control function $u(t)$;
- the **cost function** is a pay-off function where $x(t)$ is the solution for the dynamical constraint, h is the running cost and g is the terminal cost.

An IVP-ODE is defined by

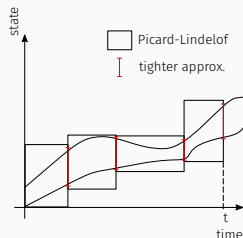
$$\begin{cases} \dot{x} = f(x) \\ x(0) \in \mathcal{X}_0 \subseteq \mathbb{R}^n, t \in [0, t_{\text{end}}] . \end{cases}$$

The goal is to compute $x(t; \mathcal{X}_0) = \{x(t; x_0) \mid x_0 \in \mathcal{X}_0\}$.

Phase 1 a priori enclosure $[\tilde{x}_j]$
of

$$\{x(t_k; x_i) \mid t_k \in [t_i, t_{i+1}], x_i \in [X_i]\}$$

Phase 2 tight enclosure of
 $[x_{i+1}]$ at time t_{i+1} .



Dynibex

- C++ library using ibex (constraint processing over real numbers);
- proof of existence and uniqueness of solution for ODEs and DAEs;
- combined with contractors (HC4), easy to use in branching algorithms;
- verification of temporal constraints.

Example of temporal constraints

- Stayed in \mathcal{A} until $\tilde{t} < t_{\text{end}}$:

$$\forall t \in [0, \tilde{t}], \{y(t; y_0) \mid y_0 \in [y_0]\} \subseteq \text{int}(\mathcal{A})$$

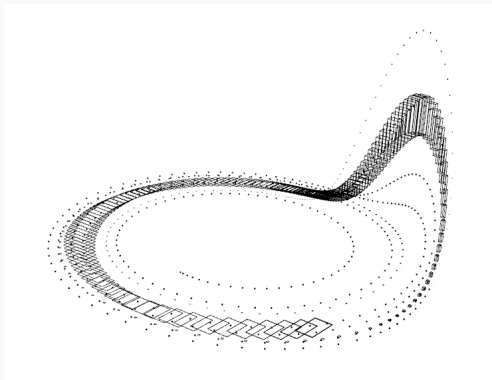
- Included in \mathcal{A} inside $[0, t_{\text{end}}]$:

$$\exists t \in [0, t_{\text{end}}], \{y(t; y_0) \mid y_0 \in [y_0]\} \subseteq \text{int}(\mathcal{A}).$$

EXAMPLE

System of Rossler: Initial states: $(0; -10.3; 0.03)$, some parameters:
 $a = 0.2, b = 0.2, c = 5.7$

$$\begin{cases} \dot{x} = -y - z \\ \dot{y} = x + ay \\ \dot{z} = b + z(x - c) \end{cases}$$



For $\mathcal{U} \subseteq [u] \in, \dot{x} \in f(x(t), [u])$ and

$$\begin{aligned}
 J(u(\cdot)) &= \int_0^T h(x(t), u(t))dt + g(x(T)) \\
 &= \sum_{i=0}^n \int_{t_i}^{t_{i+1}} h(x(t), u(t))dt + g(x(T)) \\
 &\in \sum_{i=0}^n (t_{i+1} - t_i)h([\tilde{x}_i], [u]) + g([x_T])
 \end{aligned}$$

$$\left[\begin{array}{ll}
 \max_{u(\cdot)} & J(u(\cdot)) = \int_0^T h(x(t), u(t))dt + g(x(T)) \quad (\text{cost function}) \\
 \text{s. t.} & \dot{x} = f(x(t), u(t)), \quad 0 < t \leq T \quad (\text{dynamical constraint}) \\
 & x(0) = x_0, \quad h(x(t)) \in \mathcal{H}, \quad 0 < t \leq T \quad (\text{boundary conditions}) \\
 & u(t) \in \mathcal{U}, \quad \forall t \quad (\text{bounded control})
 \end{array} \right]$$

Restriction

- Particular kind of dynamics:
 - the integral is provided by the **dynamical constraints**,
 - the set of possible control $u(t)$ is known and is discrete;
- the **cost function** is monotonic;
- the **boundary conditions** only occurs at a specific time τ .

n -mode hybrid system

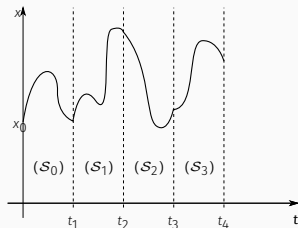
$$(\mathcal{S}_i) \begin{cases} \dot{x} = f_i(x) \\ x(t_i) = x_i \end{cases} \quad \text{in the time interval } [t_i, t_{i+1}]$$

- $f_i : \mathbb{R}^m \rightarrow \mathbb{R}^m$;
- $x_i \in \mathbb{R}^m$ is the initial condition for all modes $0 \leq i \leq n - 1$.

A sequence $\{(\mathcal{S}_1), \dots, (\mathcal{S}_k)\}$ corresponds to the switching of control law.

- x_0 is fixed;
- x_i is taken as the solution at time t_i of (\mathcal{S}_{i-1}) .

\Rightarrow not necessarily continuously differentiable.



Our problem can be modeled using the following optimization problem

$$\left[\begin{array}{ll} \max_{t_1, \dots, t_{n-1}} & g(x(\tau)) \quad \text{(cost function)} \\ \text{s. t.} & \forall 0 \leq i \leq n-1, (\mathcal{S}_i) \quad \text{(dynamical constraint)} \\ & h(x(\tau)) > 0 \quad \text{(reachability constraint)} \\ & \tau \in [t_{n-1}, t_n] \end{array} \right]$$

with

- the **decision variables** $t_1, \dots, t_{n-1} \in \mathbb{R}_+^n$ the search space for the different times;
- the **cost function** $g : \mathbb{R}^m \rightarrow \mathbb{R}$ on the state variable at given time $\tau \in [t_{n-1}, t_n]$;
- some constraints defined by the **dynamical systems** (\mathcal{S}_i) and the times t_i ;
- a **reachability constraint** using $h : \mathbb{R}^m \rightarrow \mathbb{R}$.

EXAMPLE: GODDARD'S ROCKET

Model of the ascent of a rocket in the atmosphere:

with the parameters

$$\left[\begin{array}{l} \max \quad m(T) \\ \text{s.t.} \quad \dot{r} = v \\ \dot{v} = \frac{u - Av \exp(k(1-r))}{m} - \frac{1}{v^2} \\ \dot{m} = -bu \\ u(\cdot) \in [0, 1] \\ r(0) = 1, v(0) = 0, m(0) = 1 \\ r(T) \geq \mathcal{R}_T \end{array} \right] \begin{array}{l} \cdot b = 2, \\ \cdot T_{\max} = 0.2, \\ \cdot A = 310, \\ \cdot k = 500, \\ \cdot r_0 = 1, v_0 = 0, m_0 = 1, \\ \cdot \mathcal{R}_T = 1.01. \end{array}$$

According to the time t :

$$u(t) = \begin{cases} 3.5 & \text{for } t \in [0, t_1] \quad (\mathcal{S}_0) \\ 3.5 \tanh(1+t) & \text{for } t \in [t_1, t_2] \quad (\mathcal{S}_1) \\ 0 & \text{for } t \in [t_2, T] \quad (\mathcal{S}_2) \end{cases}$$

Algorithm 1: $\text{simu}(t_1, t_2, \text{max})$ – simulates the system from 0 to T .

Input: time t_1, t_2 to switch dynamics; current maximum mass max

Output: the mass m or 0 if simulation will not produce a better solution

$([r_{t_1}], [v_{t_1}], [m_{t_1}]) \leftarrow$ simulation of x_0 using (\mathcal{S}_0) from 0 to t_1 ;

if $\bar{m}_{t_1} \leq \text{max}$ then

\perp return 0;

$([r_{t_2}], [v_{t_2}], [m_{t_2}]) \leftarrow$ simulation using (\mathcal{S}_1) from t_1 to t_2 ;

if $\bar{m}_{t_2} \leq \text{max}$ then

\perp return 0;

$([r_T], [v_T], [m_T]) \leftarrow$ simulation using (\mathcal{S}_2) from t_2 to T ;

if $r_T \geq \mathcal{R}_T$ then

 return $[m_T]$;

else

\perp return 0;

Algorithm 2: finds the optimal switching times

Input: set of dynamics $\{(\mathcal{S}_0), (\mathcal{S}_1), (\mathcal{S}_2)\}$

Output: switching times $t_{1,max}$ and $t_{2,max}$

$max \leftarrow 0;$

for $t_1 \leftarrow 0$ to $T - \epsilon$ do

 for $t_2 \leftarrow t_1 + \epsilon$ to T do

$[\underline{m}, \overline{m}] \leftarrow \text{simu}(t_1, t_2);$

 if $\underline{m} > max$ then

$\underline{m} \leftarrow max;$

$t_{1,max} \leftarrow t_1;$

$t_{2,max} \leftarrow t_2;$

 else

 break;

 // Due to monotonicity of the cost

 function

return $t_{1,max}$ and $t_{2,max};$

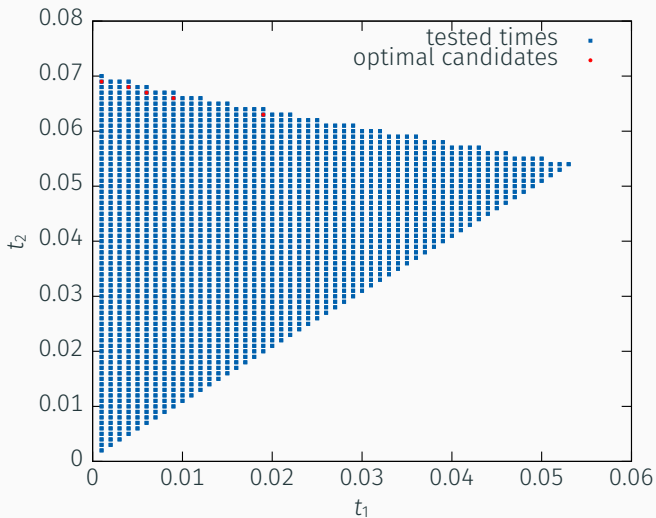


Figure 1: Mesh for t_2 w.r.t. t_1

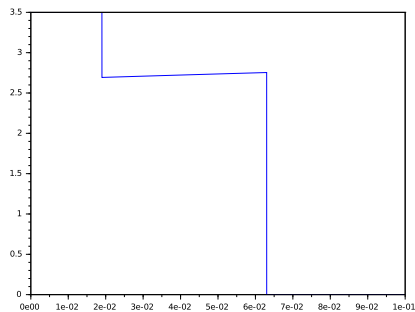


Figure 2: Optimal controller

Conclusion

- Promising results on the computation of optimal switching mode;
- easy-to-use tool development;
- benefits shown on an example.

Perspectives

To release restrictions to handle the problem in its generality



Robert H Goddard.

A method of reaching extreme altitudes.

Nature, 105:809–811, 1920.



Knut Graichen and Nicolas Petit.

Solving the goddard problem with thrust and dynamic pressure constraints using saturation functions.

IFAC Proceedings Volumes, 41(2):14301–14306, 2008.